



Preserving Data Integrity in Peer-to-Peer Systems

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What is Peer-to-Peer?

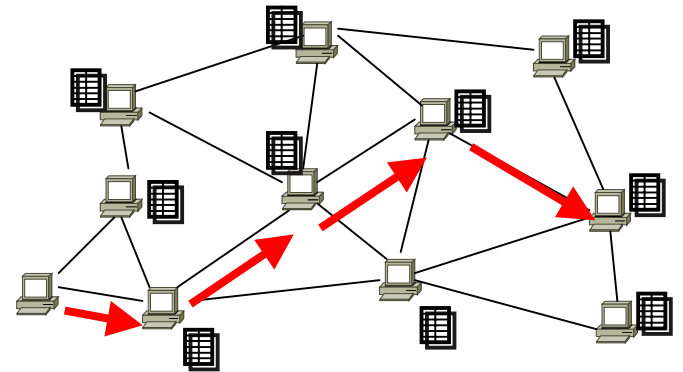
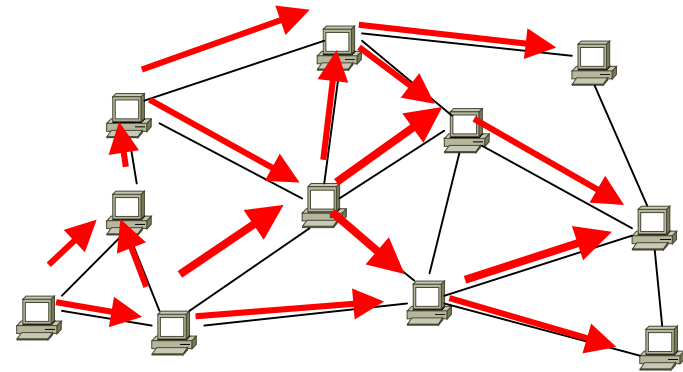
- Peer-to-Peer (P2P) concept:
 - Leverage idle resources
- Definition:
 - Self-organizing distributed system
 - Nodes provide and receive services in cooperative effort
- Features:
 - Scalability
 - Availability
 - Fault tolerance, etc.

Search in P2P

- A key operation
- Search query:
 - given name or keyword attributes of content, where is it?
- Search response:
 - a set of index entries pointing to replica nodes storing the content
- Index entry:
 - (key, value) pair
 - Key = name of content
 - Value = IP address of serving peer

Great for (illegal) file-sharing!!

- Unstructured
 - Query flooding
 - Gnutella, FreeNet
- Structured
 - Single query path
 - CAN, Chord, Pastry, Tapestry
- Anything else?



Problem Characteristics

- Participating entities are
 - Autonomous
 - Mutually distrustful
 - Mutually dependent

Example: Digital Preservation of on-line published material

Traditional Library Model

- Goal: Preserve access to important documents for posterity
- On behalf of their institution, libraries
 - Acquire and distribute lots of paper copies of important materials
 - Give access to local readers
 - Lend copies to other libraries
- It is hard to destroy all copies

Transition to Digital Media

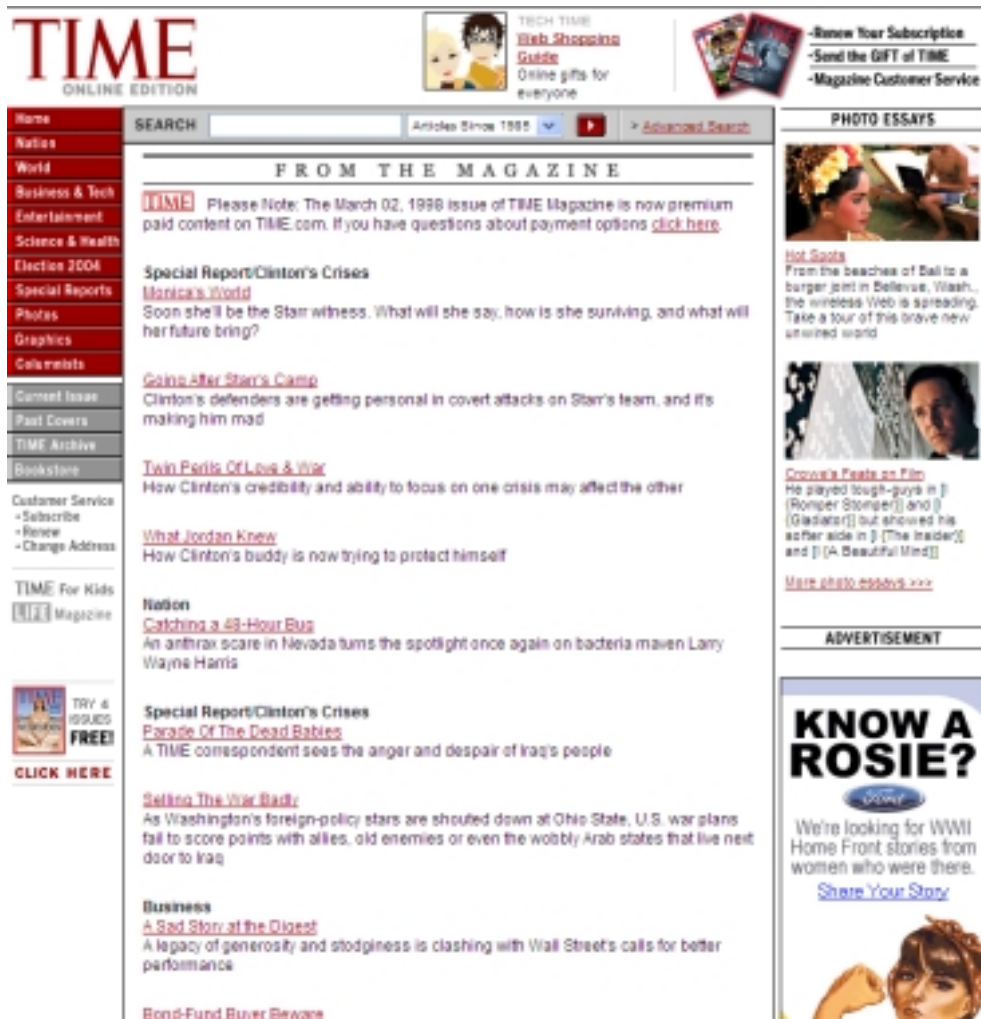
- Resources: electronic journals, proceedings, etc.
 - Publisher rents access to materials
- Problem:
 - Libraries no longer own the bits
 - Libraries vulnerable to disappearance, failure, misbehavior of publishers

Example: Time Magazine



Essay by Bush Sr. published in paper-version of March 2, 1998 issue

Online Version Removed



- Online version has disappeared
- Online table of contents modified
- It's as if article never existed in this issue!

LOCKSS Goals

- Lots Of Copies, Keep Stuff Safe
- Emulate traditional model for on-line publishing
- Make it easy for libraries to
 - Own, rather than lease, materials
 - Preserve and provide access to local patrons
- Make it easy for publishers to
 - Provide content for preservation and archiving
 - Without the responsibility for perpetual presence
 - With minimal risk to their business model

LOCKSS Approach

- Build p2p community of libraries
- Audit and repair their contents with
 - No centralized control (Autonomous)
 - Mutual distrust
 - Very low-cost hardware, operation and administration (Mutually Dependent)
 - A long-term horizon; I.e., preserve for decades
- Must anticipate natural bit degradation
- Must anticipate sustained attacks

Opinion Polls

- Obtaining full consensus is difficult
- Each peer holds
 - Reference list of peers it has discovered
 - Friends list of peers it knows externally
- Periodically (faster than rate of bit rot)
 - Takes a sample of the reference list
 - Invites the chosen peers to send a hash of their copy of the document

Opinion Polls (cont'd)

- Peer compares votes with local copy
- If landslide agreement, the peer is happy
- If landslide disagreement, the peer repairs
 - To repair, the peer gets the copy of somebody who disagreed and then reevaluates the same votes
- If poll is inconclusive, the peer raises alarm
 - Alarms are built-in intrusion detection

Reference List Update

- Take out voters in the poll
 - So that the next poll is based on different group
- Replenish with some “strangers” and some “friends”
 - Strangers: Accepted nominees proposed by voters
 - Friends: From the friends list
 - The measure of favoring friends is called *churn factor*

Adversary Goals

- Top adversary goal: **Stealth Modification**
 - Modify documents unobtrusively
 - Hard to reinstate original content after large proportion of peers have had their documents modified
- Other goals
 - Slow the system down
 - Discredit the system
 - Obtain benefits without contributing
 - Obtain content illicitly

LOCKSS Defenses

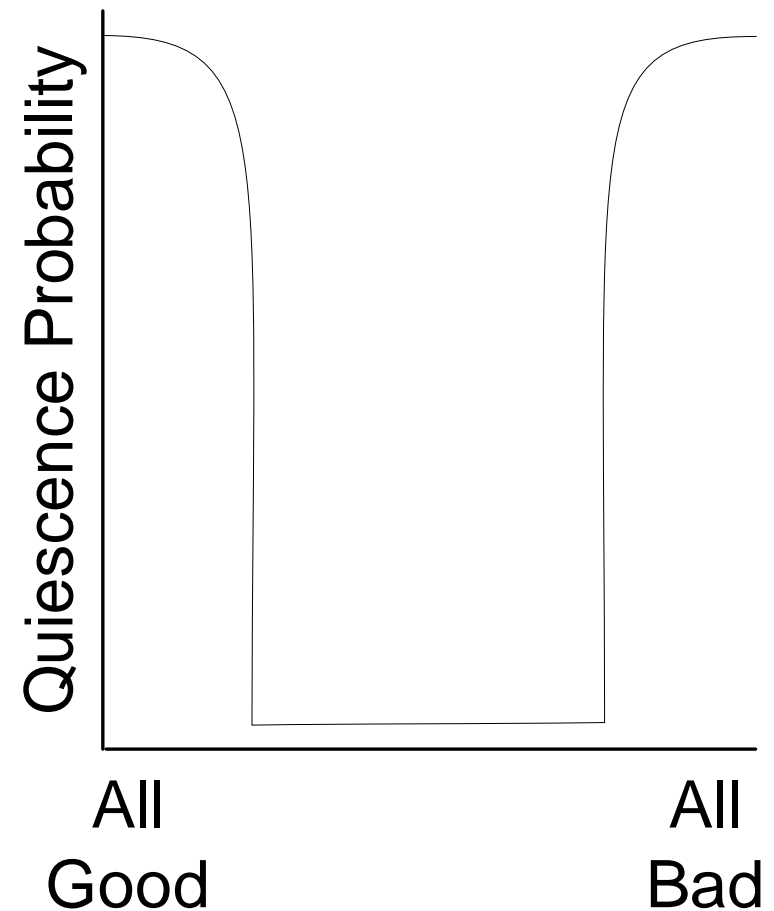
- Limit the rate of operation
- Bimodal system behavior
- Churn friends into reference list

Limit the rate of operation

- Peers determine their rate of operation autonomously
 - Adversary must wait for the next poll to attack through the protocol
- No operational path is faster than others
 - Artificially inflate “cost” of cheap operations
 - No attack can occur faster than normal ops

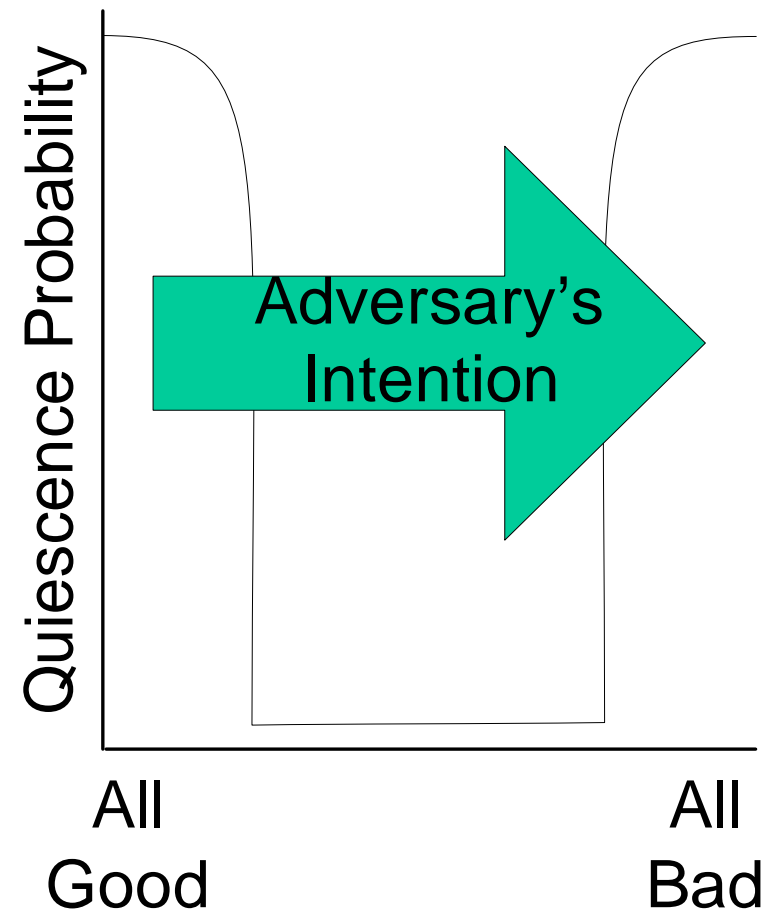
Bimodal System Behavior

- When most replicas are the same, no alarms
- In between, many alarms
- To get from mostly correct to mostly wrong replicas, system must pass through “moat” of alarming states



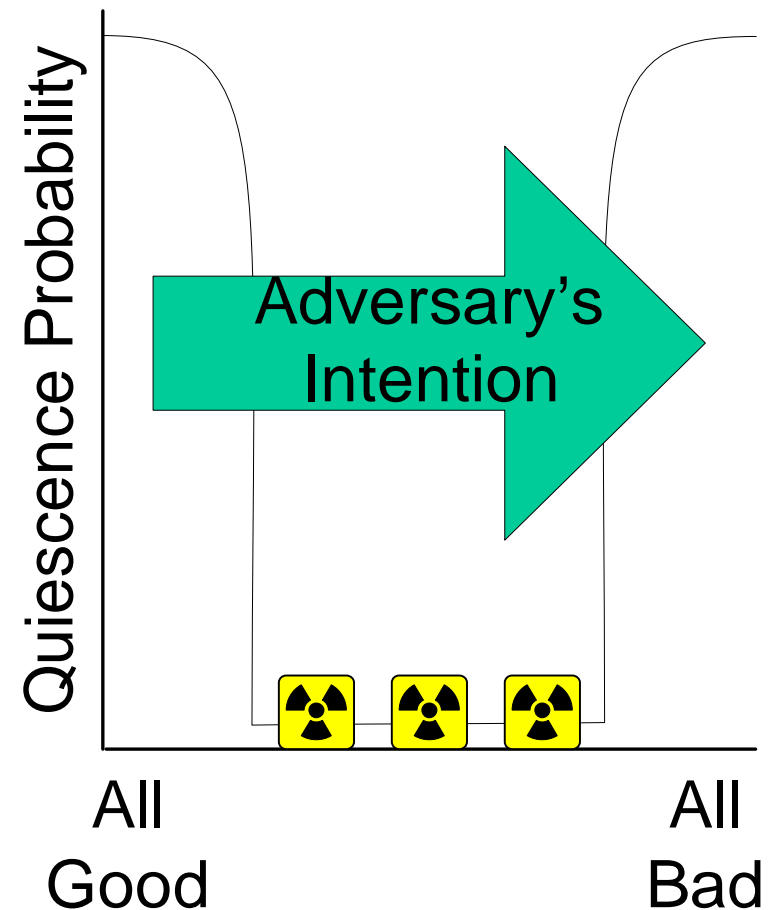
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Churn Friends into Reference List

- Churn adjusts the bias in the reference list
- High churn favors friends
 - Reduces the effects of Sybil attacks
 - But offers easy targets for focused attack
- Low churn favors strangers
 - It offers Sybil attacks free reign
 - Bad peers nominate bad; good peers nominate some bad
 - Makes focused attack harder, since adversary can predict less of the poll sample
- Goal: strike a balance

Evaluation Methodology

- Model a very powerful, realistic adversary
- Identify major goals of adversary attacks
- Devise and implement rational strategies
- Measure the impact of each strategy
 - locally (on library patrons)
 - globally (on document survival)

Adversary Model

- Unlimited identities
 - Purchased (cheap) or spoofed (cheaper)
- Exploits common implementation bugs
 - Exploited peer is *subverted*
- Perfect coordination
 - Instantaneous communication with and control of subverted peers
 - Load balancing of attack effort
 - Flawless content preservation

Stealth Modification Strategy

- A peer's reference list affects outcomes of polls it will call
- The stealth adversary
 - First, quietly gains a strong foothold in the reference list of a peer
 - Then, attacks when a poll will be landslide win in favor of adversary's copy
 - Must consistently win polls to succeed

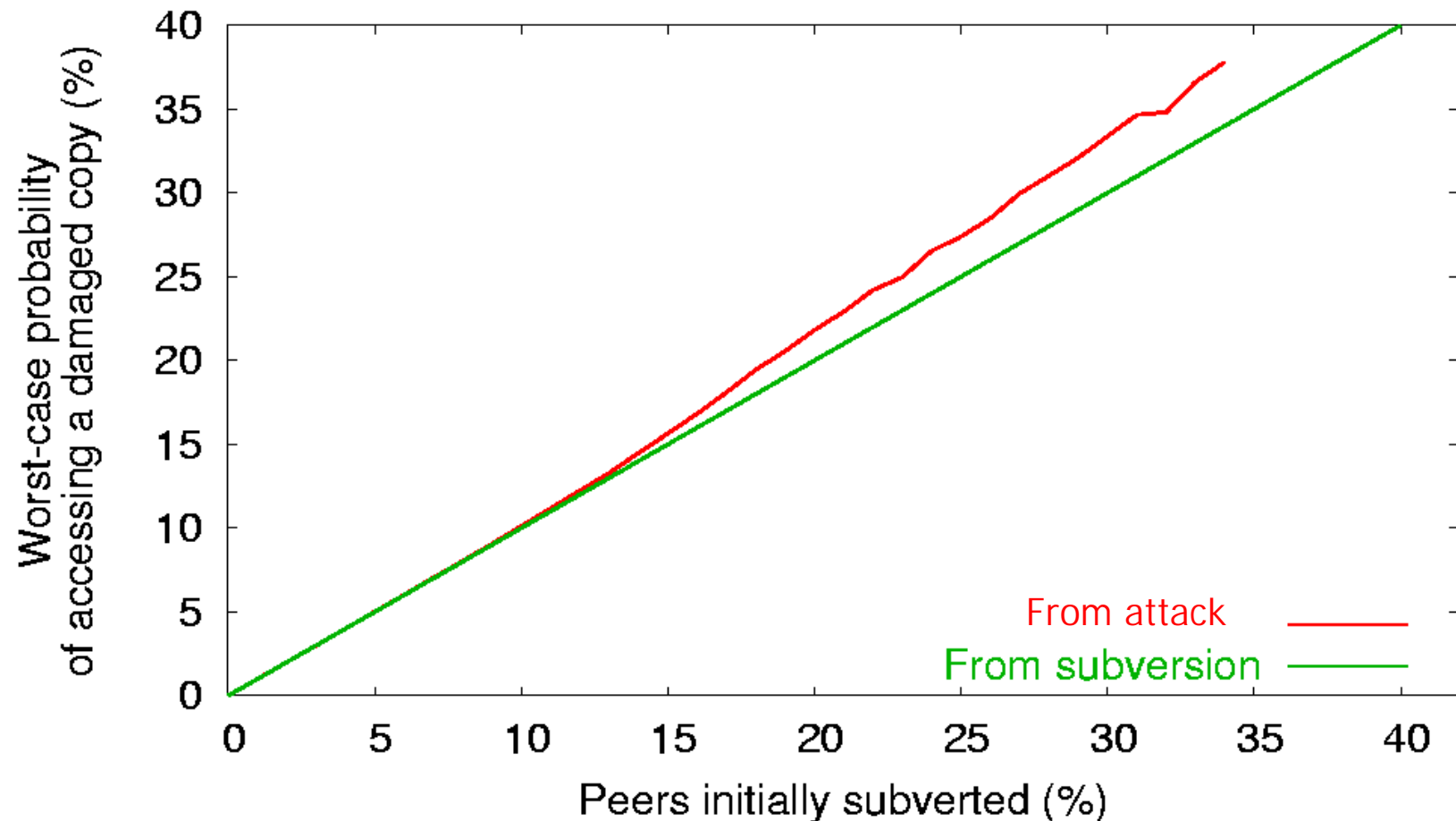
Evaluation

- We use Narses, an application-layer protocol simulator
- Scenarios
 - 1000 original peers, in clusters of friends
 - Initially, 0 – 40% are subverted
 - Lurk for up to 20 years
 - Attack for up to 10 more years
 - Report worst-cases over ~200 runs per data point (recent results)

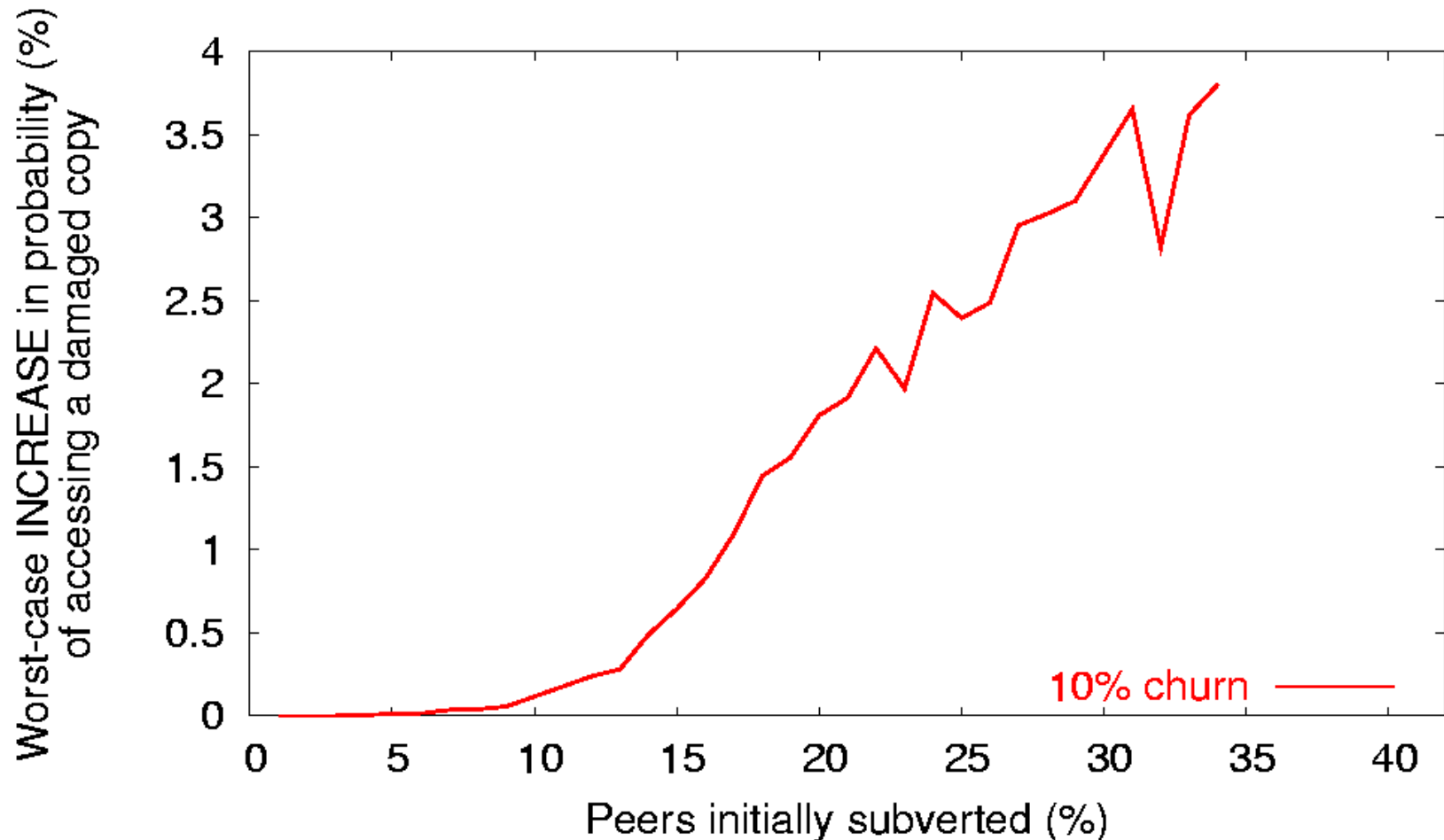
Metrics

- Metrics
 - What's the probability that an access reaches a bad replica
 - What's the probability that the document is damaged irrecoverably
- **How big is the effect of the worst protocol attack on top of the effect of the initial subversion?**

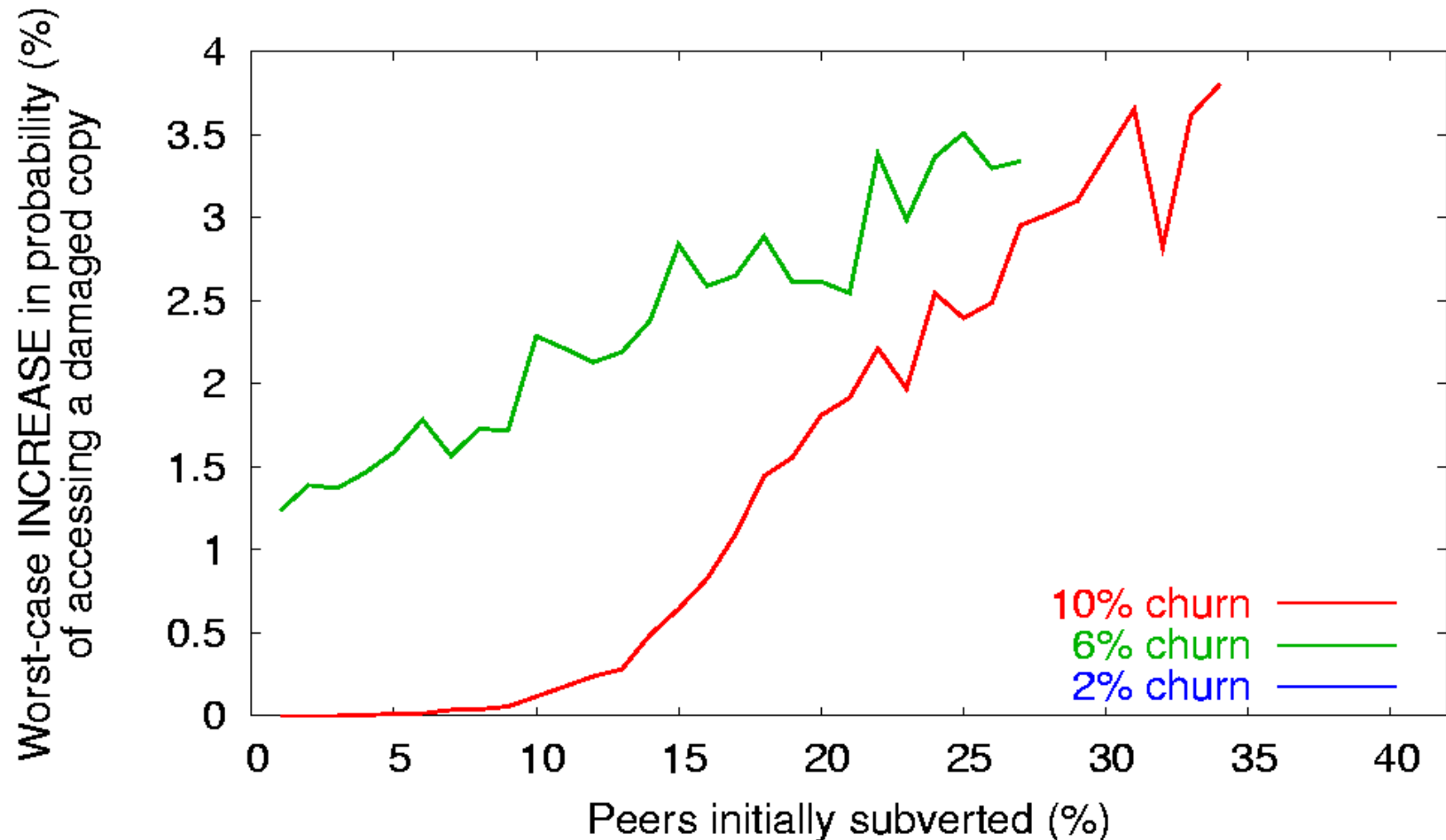
Probability of Accessing Bad Copy



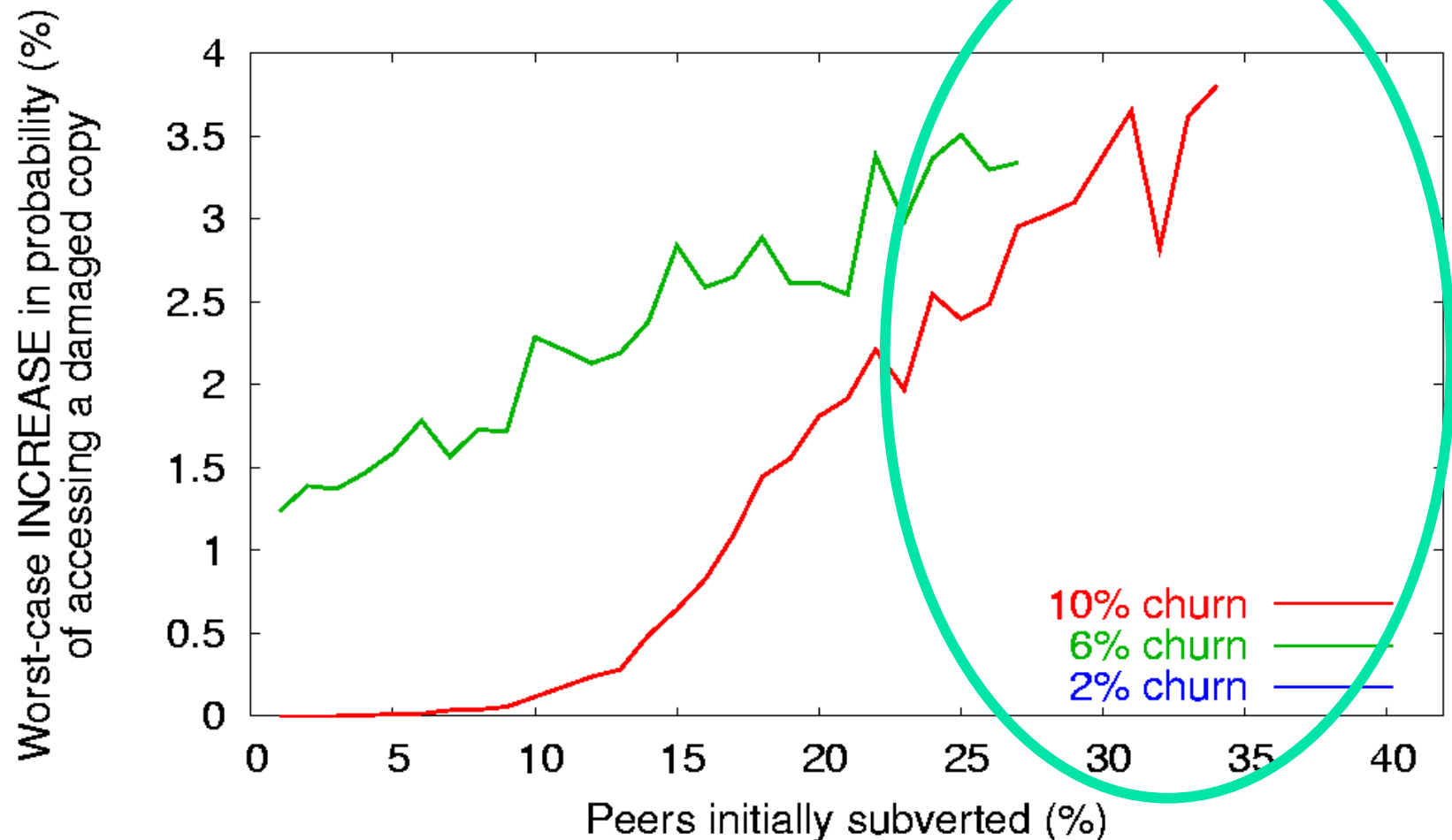
Probability of Accessing Bad Replica (Incremental)



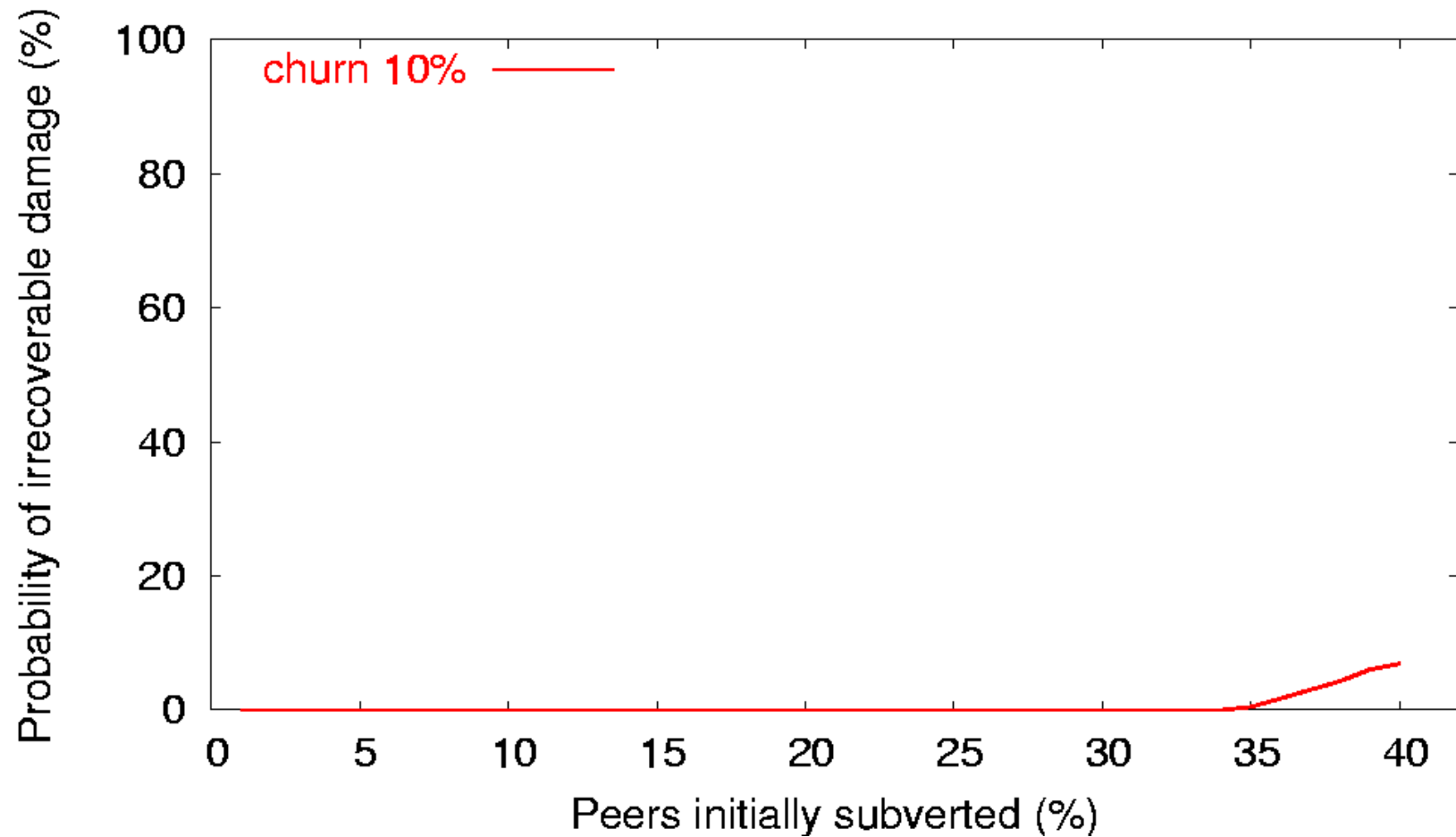
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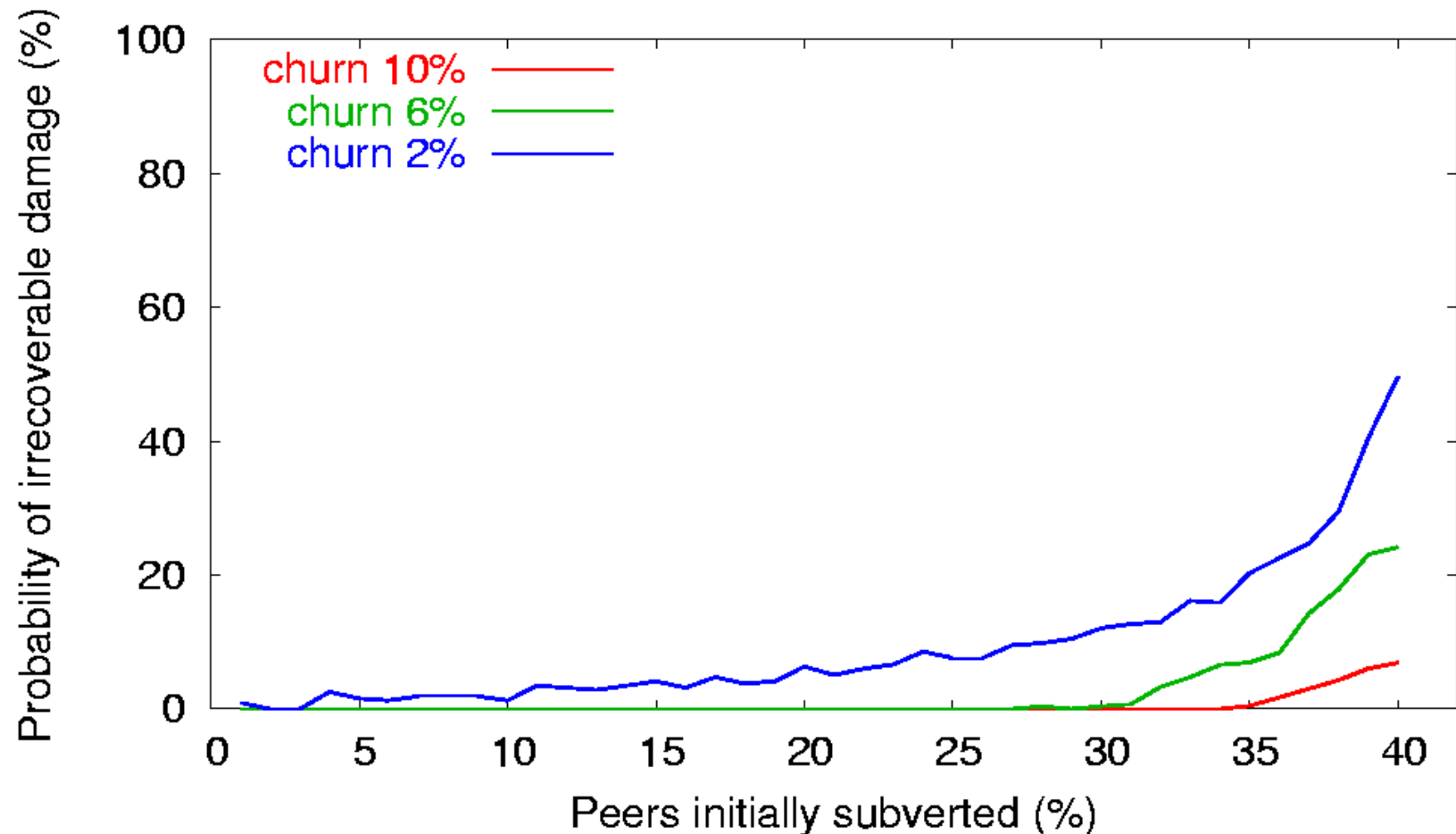
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Probability of Irrecoverable Damage



Probability of Irrecoverable Damage



Alternatives

- Use super-fabulous RAID
 - Can be complementary, but alone cannot ensure survivability when failures do occur (e.g., human error)
- Encrypt or sign to ensure integrity
 - Preserving public keys just as hard a problem
- Boost efficiency with erasure codes etc.
 - Storage space is not an issue
 - All replicas must be whole

Next Steps

- Explore the parameter space
 - What quorum sizes are necessary?
 - Frequency of polls vs. rate of undetected medium faults vs. probability of adversary success
- Enlarge bestiary of attackers
 - Attrition attacks (e.g., DDoS)
 - Hybrid attacks (e.g., stealth modification during DDoS weakening)
- Expand to other application domains

Applications

- Academic journals
 - Append-only updates
- Government documents
 - Large number, frequent updates
- High-resolution scans of artwork
 - “Rare-bits”
- Scientific data
 - Large volumes (terabytes) of data

Conclusions

- P2P is more than file-sharing
- P2P good for applications with:
 - Autonomous entities
 - Mutually distrustful entities
 - Mutually dependent entities
- One example: LOCKSS, a P2P digital preservation system

Status

- Results for stealth adversary
 - Resistant to attacks for low subversions
 - Degrades gracefully for greater subversions
- Status
 - Promising results for other attacks (DDOS)
 - To be deployed at ~100 libraries across the globe in 2004
 - For more info:

<http://www.eecs.harvard.edu/~mema/>